

Non-Research Tips for Information Science Researchers (Summer 2026)

Apr 22, 2026

Week 2: Equations and pseudo-codes

<https://non-research-tips.github.io/2026>



Yusuke Matsui
(UTokyo)

Schedule

| Date (2026) | Contents | Presented by |
|---------------------------|---|---|
| Week 1, Apr 15 | Introduction. Review of fundamental concepts | Yusuke, Koya, Yuki, Jun |
| Week 2, Apr 22 | Equations and pseudo-codes | Yusuke Matsui |
| Week 3, May 13 | Slides | Koya Narumi |
| Week 4, May 20 | Research community | Jun Kato |
| Week 5, May 27 | Figures | Koya Narumi |
| Week 6, June 1 | 3DCG illustrations | Yuki Koyama |
| Week 7, June 10 | Invited Talk 1: Research tips in the private sector: From topic selection to social implementation (temp.) | Dr. Antonio Tejero de Pablos (CyberAgent) |
| Week 8, June 17 | Tables and plots | Yusuke Matsui |
| Week 9, June 24 | Invited Talk 2: Portable and Reproducible Research Environments in the Age of AI Agents | Dr. Mai Nishimura (OSX) |
| Week 10, July 1 | DevOps for research | Jun Kato |
| Week 11, July 8 | Videos | Koya Narumi |
| Week 12, July 15 | Final presentations | YOU |



Author



The data x is...

It's obvious from the context what x is.

Author



The data x is...

Reviewer 2



It's obvious from the context what x is.

What is x ? A scalar? A vector?
Hard to read... Reject!

Author



The data x is...

Reviewer 2



It's obvious from the context what x is.

What is x ? A scalar? A vector?
Hard to read... Reject!



The data $x \in \mathbb{R}^D$ is...



Let me explicitly define x as a D -dim vector.

Ok, x is a D -dim real-valued vector. This author uses a bold alphabet as a vector. Got it. Keep reading...

Author



The data x is...

Reviewer 2



Writing a clear, simple, and precise mathematical description is extremely important!



The data $\mathbf{x} \in \mathbb{R}^D$ is...



Let me explicitly define \mathbf{x} as a D -dim vector.

Ok, \mathbf{x} is a D -dim real-valued vector. This author uses a bold alphabet as a vector. Got it. Keep reading...

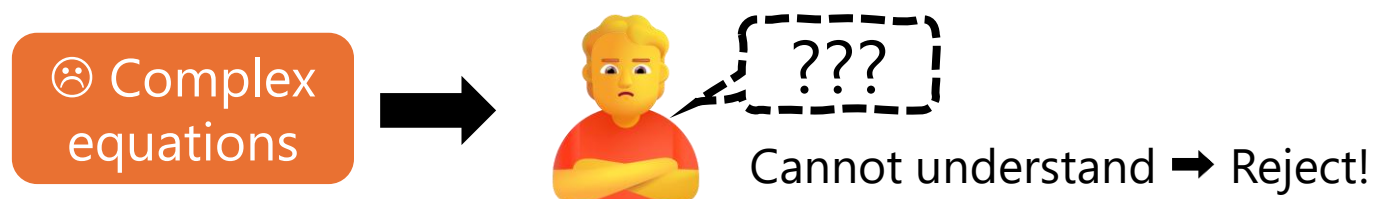
Equations and pseudo-codes

Introduce the guideline of writing equations. Three important things:

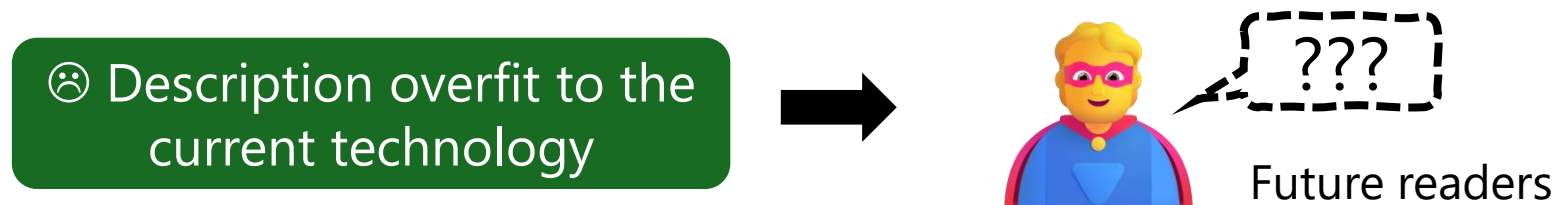
1. Describe **precisely what you want to say** without making mistakes.



2. Be **sufficient**, not overly complex or ambiguous.



3. Be understandable even if it is read **10 years later**.



Equations and pseudo-codes

- Main target audience:
 - ✓ Researchers in CV, NLP, or related fields.
- Different fields have very different conventions.
 - ✓ If the content in this lecture differs from your field's convention, **please always follow your field's convention.**
 - ✓ e.g., I suggest using a bold font for vectors, but your field might never use bold.



"OK" in US

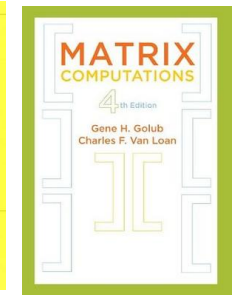
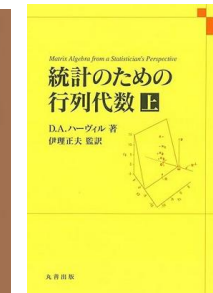
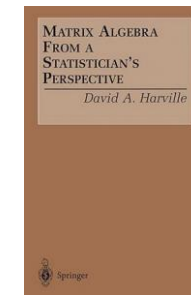


"OK" in Japan

Reference

Dictionary for notation

- D. A. Harville, "Matrix Algebra From a Statistician's Perspective", Springer, 2000.
- D. A. ハーヴィル, "統計のための行列代数上・下", 丸善出版, 2012.
- G. H. Golub and C. F. Van Loan, "Matrix Computations, 4th edition", Johns Hopkins University Press, 2012.



Very precise description for Transformer

- M. Phuong and M. Hutter, "Formal Algorithms for Transformers", arXiv 2022. <https://arxiv.org/abs/2207.09238>

Higher-order tensor

- T. G. Kolda and B. W. Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.
- 横田達也, "テンソル分解の基礎と応用", MIRU チュートリアル, 2022. <https://speakerdeck.com/yokotatsuya/tensorufen-jie-falseji-chu-toying-yong-miru2022tiyutoriariu>
- 横田達也, "テンソルデータ解析の基礎と応用", コロナ社, 2024



The original document for this lecture

- 松井勇佑, "工学系の卒論生のための数式記述入門", GitHub, 2021. https://github.com/mti-lab/math_writing

工学系の卒論生のための数式記述入門

東京大学 情報理工学系研究科 講師 松井勇佑

初版 2021年1月7日 / 最終更新 2023年4月26日

- 本資料の GitHub リポジトリ: https://github.com/mti-lab/math_writing
- 著者のウェブページ: <http://yusukematsui.me>

1 はじめに

本資料は、初めて技術文章を書く学生向けの、数式記述の入門書です。自分のアイデアを式として記述する作業は難しく、最初のうちはうまく書けません。本資料は、そんな学生向けに、数式をどう記述すればいいかのガイドラインを示します。特に、よくありがちな、変数がゴチャゴチャしてわかりにくくなったり、アルファベットの

しょう。

- 東京大学松尾先生による記事: 読み物としても面白い、松尾先生による論文執筆指南です。短い文章で重要な事項が凝縮されており、必読です。
 - 松尾くみの論文の書き方
 - 松尾くみの論文の書き方: 英語論文
- D. A. ハーヴィル, 統計のための行列代数 上下: これは線形

次に、数式の表記 (notation) に関して参考になる資料を紹介いたします。本記事では、これらを参考に、松井の考える notation のベストプラクティスを紹介します。

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

Basic notation for variables

- Having a good notation for variables is extremely important.
- The most important rule: **to be consistent**
 - ✓ Bold for vectors (recommended):
 - This is $\mathbf{x} \in \mathbb{R}^3 \dots$ Here, $\mathbf{y} \in \mathbb{R}^3 \dots$
 - ✓ Non-bold for vectors (ok):
 - This is $x \in \mathbb{R}^3 \dots$ Here, $y \in \mathbb{R}^3 \dots$
 - ✓ Mixed (**REALLY BAD!** 🤪 🤪 🤪 🤪):
 - This is $\mathbf{x} \in \mathbb{R}^3 \dots$ Here, $y \in \mathbb{R}^3 \dots$
- If you decide your way, keep following the way **throughout the paper.**

Basic notation for variables

- For all variables, explicitly write a **domain**:

\in: "an element of"

domain: a set, e.g., \mathbb{R}

$$x \in \mathbb{R}$$

x is an element of real numbers
→ x is a real number

- Including a domain is crucial; much easier to understand.
- If you think you can skip a domain, you're likely incorrect by 90%.
- Writing a domain only requires a bit more space. **No side effects.**

- An exceptional case is when you hide a domain **intentionally**.

- ✓ Avoid unnecessarily complex descriptions. (I'll explain later)

- ✓ e.g., a feature volume for time series $V \in \mathbb{R}^{H \times W \times T}$ 🤔 T depends on $V...$

First things first: real numbers, natural numbers, etc

➤ **Blackboard bold** is used for special symbols such as real numbers.

✓ TeX: `\mathbb`

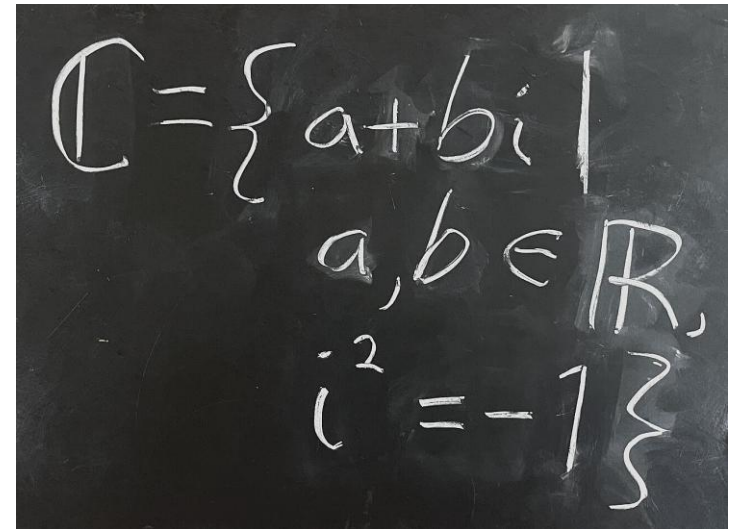
✓ Powerpoint: e.g., `\doubleR`

➤ Examples:

✓ \mathbb{R} : Real numbers

✓ \mathbb{N} : Natural numbers

✓ \mathbb{Z} : Integers



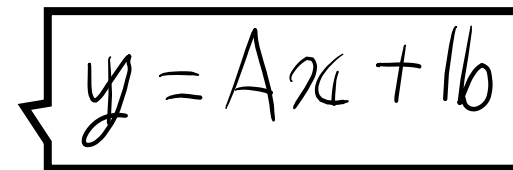
A photograph of a chalkboard with white chalk writing. The text reads: $\mathbb{C} = \{ a+bi \mid a, b \in \mathbb{R}, i^2 = -1 \}$. The set notation is written with a large curly brace, and the conditions are written on separate lines.

https://en.wikipedia.org/wiki/Blackboard_bold

➤ Blackboard bold is typically used only in specific contexts.

✓ 🚫 Don't recommend using it for a vector: $\mathbb{x} \in \mathbb{R}^3$

✓ (although we do so when writing a vector by hand...)



A handwritten equation $y = Ax + b$ enclosed in a black rectangular box. The text is written in a cursive, hand-drawn style.

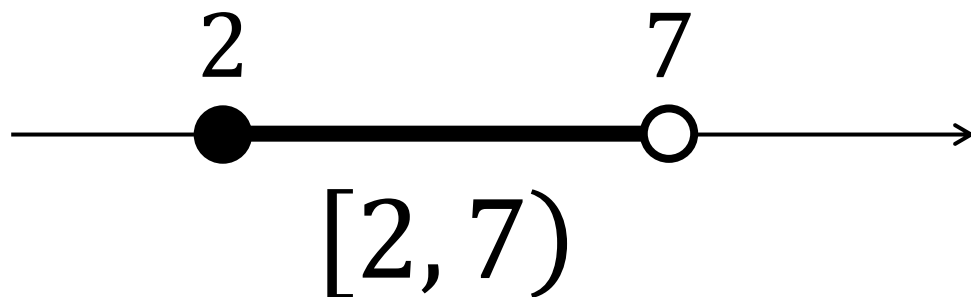
Summary

| Type | How to write | Example of domain | Example of values |
|-------------------|--|--|--|
| Scalar | Lowercase or uppercase | $a \in \mathbb{R}. b \in \mathbb{N}. K \in \{10, 20, 30\}.$ | $a = 3.2. b = 13. K = 20.$ |
| Vector | Lowercase and bold ➤ TeX: <code>\mathbf</code> or <code>\bm</code> ➤ Powerpoint: bold | $\mathbf{x} \in \mathbb{R}^3. \mathbf{b} \in \{0, 1\}^B.$ | $\mathbf{x} = [0.1, 0.2, 0.3]^\top. \mathbf{b} = [0, 1, 1, 0]^\top.$ |
| Matrix and Tensor | Uppercase and bold ➤ TeX: <code>\mathbf</code> or <code>\bm</code> ➤ Powerpoint: bold | $\mathbf{A} \in \mathbb{R}^{2 \times 3}. \mathbf{I} \in [0, 1]^{H \times W \times 3}.$ | $\mathbf{A} = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 4 \end{bmatrix}.$ |
| Set | Calligraphy font in uppercase ➤ TeX: <code>\mathcal</code> ➤ Powerpoint: e.g., <code>\scriptS</code> | $\mathcal{S} \subset \mathbb{N}.$ | $\mathcal{S} = \{2, 4, 8, 16\}.$ |

Scalar

- Recommend: Lowercase or uppercase (e.g., $a \in \mathbb{R}$, $K \in \mathbb{N}$, $\mu \in \mathbb{C}$)
- If you can specify the range tightly, it's more informative.

| Example | Meaning |
|-------------------------|--|
| $a \in [2, 7]$ | $2 \leq a \leq 7$ More informative than $a \in \mathbb{R}$ |
| $a \in (2, 7)$ | $2 < a < 7$ |
| $a \in [2, 7)$ | $2 \leq a < 7$ |
| $a \in \{2, 7\}$ | $a = 2$ or $a = 7$ |
| $a \in \{2, \dots, 7\}$ | If naturally interpreted, $a = 2$ or $a = 3$ or ... or $a = 7$. |



Range (remember your high-school math!)

All brackets are different and have various meanings

() : Parentheses

[] : Square brackets

{ } : Curly brackets

Vector

- Recommend: Lowercase and bold, e.g.,
 - ✓ $\mathbf{x} \in \mathbb{R}^3$ 3-dimensional real-valued vector
 - ✓ $\mathbf{b} \in \{0, 1\}^B$ B -dimensional binary vector
- How to write
 - ✓ TeX: `\mathbf` or `\bm`
 - ✓ Powerpoint: set bold
- `\mathbf` or `\bm`?
 - ✓ Depends on the TeX style.
 - ✓ Render both, compare them, and select the best one.

| | |
|----------------------|---|
| Default | $x, a, 1, \mu$ |
| <code>\mathbf</code> | $\mathbf{x}, \mathbf{a}, \mathbf{1}, \mu$ |
| <code>\bm</code> | $\mathbf{x}, \mathbf{a}, \mathbf{1}, \mu$ |

Vector

- Recommend: Lowercase and bold, e.g.,
 - ✓ $x \in \mathbb{R}^3$ 3-dimensional real-valued vector
 - ✓ $b \in \{0, 1\}^B$ B -dimensional binary vector
- How to write
 - ✓ TeX: `\mathbf` or `\bm`
 - ✓ Powerpoint: set bold
- `\mathbf` or `\bm`?
 - ✓ Depends on the TeX style.
 - ✓ Render both, compare them, and select the best one.

Default

$x, a, 1, \mu$

Non-italic. Very different from the default.

`\mathbf`

$\mathbf{x}, \mathbf{a}, \mathbf{1}, \mu$

Special characters may not work

`\bm`

$\mathbf{x}, \mathbf{a}, \mathbf{1}, \mu$

Italic. Natural but hard to distinguish from the default

Vector

➤ Row-vector or column-vector?

- ✓ Recommend: If you write a vector (e.g., $\mathbf{x} \in \mathbb{R}^3$), assume that all vectors are column-vectors.

$$\mathbf{x} = [1, 2, 3]^T = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix} \quad \text{Column-vector. OK.}$$

$$\mathbf{y} = [1, 2, 3] \quad \text{Row-vector. Only write this if you explicitly need it.}$$

➤ Why should we be careful?

- ✓ Need to be careful when performing matrix operations:

$$\text{Given } \mathbf{A} \in \mathbb{R}^{3 \times 3}, \begin{cases} \mathbf{Ax} \text{ can be computed} \\ \mathbf{Ay} \text{ cannot be computed } (\mathbf{Ay}^T \text{ can be computed}) \end{cases}$$

Vector

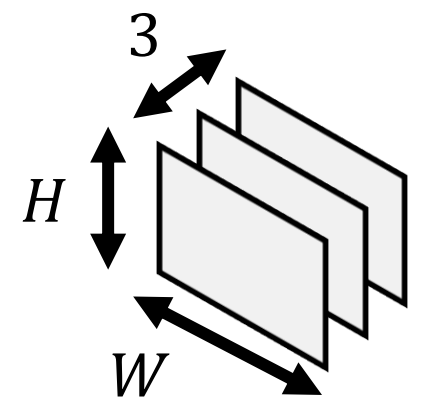
- More strictly, just writing " \mathbb{R}^D " cannot decide it's column or row.
 - ✓ Explicitly saying D -dim column-vectors: $\mathbf{a} \in \mathbb{R}^{D \times 1}$
 - ✓ Explicitly saying D -dim row-vectors: $\mathbf{b} \in \mathbb{R}^{1 \times D}$
- Recommend: Decide in your mind that $\mathbf{c} \in \mathbb{R}^D$ is an abbreviated notation of $\mathbf{c} \in \mathbb{R}^{D \times 1}$
- Important rule: Again, make it to be **consistent** throughout the paper.
 - ✓ Recommend: "column-vector $\mathbf{x} \in \mathbb{R}^3$ and column-vector $\mathbf{y} \in \mathbb{R}^3$ "
 - ✓ Not recommend but OK: "row-vector $\mathbf{x} \in \mathbb{R}^3$ and row-vector $\mathbf{y} \in \mathbb{R}^3$ "
 - ✓ 🙄 NO!!!: "column-vector $\mathbf{x} \in \mathbb{R}^3$ and row-vector $\mathbf{y} \in \mathbb{R}^3$ "

Recommend this way if you need a row vector

Same description but sometimes column, sometimes row??? Confusing. Reject!



Matrix and tensor



- Recommend: Uppercase and bold, e.g.,
 - ✓ $A \in \mathbb{R}^{2 \times 3}$ 2x3 real-valued matrix
 - ✓ $I \in [0, 1]^{H \times W \times 3}$ HxWx3 3rd-order tensor. Each element is in $[0, 1]$
- Bold or not?
 - ✓ Non-bold may be more usual.
 - ✓ But a non-bold uppercase alphabet can be misinterpreted as a scalar (e.g., K is a matrix? Scalar?). Thus, I prefer bold for a matrix.
- Higher-order tensor is often used in CV as an image is 3rd-order tensor.
 - ✓ For higher-order tensors, one may use `\mathsf` or bold `\mathcal`

$$A \in \mathbb{R}^{2 \times 3 \times 4}$$

$$\mathcal{A} \in \mathbb{R}^{2 \times 3 \times 4}$$

Matrix and tensor

| Example | Meaning |
|---|--|
| $a \in \mathbb{R}$ | Real-valued scalar |
| $\mathbf{a} \in \mathbb{R}^2$ | Two-dimensional real-valued vector |
| $\mathbf{A} \in \mathbb{R}^{2 \times 3}$ | 2x3 real-valued matrix |
| $\mathbf{B} \in [0, 1]^{2 \times 3}$ | 2x3 matrix, where each element is in $[0, 1]$ |
| $\mathbf{C} \in [0, 1]^{2 \times 3 \times 4}$ | 2x3x4 tensor, where each element is in $[0, 1]$ |
| $\mathbf{D} \in \{-1, 0, 1\}^{2 \times 3}$ | 2x3 matrix, where each element is either -1 or 0 or 1 |
| $\mathbf{E} \in \{0, \dots, 100\}^{2 \times 3}$ | 2x3 matrix, where each element is either 0 or ... or 100 |
| $\mathbf{f} \in \mathbb{N}^{1 \times ab}$ | (ab) -dimensional row vector of natural numbers |
| $\mathbf{F} \in \mathbb{N}^{a \times b}$ | $a \times b$ matrix of natural numbers |



Reshape

Matrix and tensor

- Accessing an element of a vector: Given $\mathbf{a} \in \mathbb{R}^3$,
 - ✓ i -th element: $a_i \in \mathbb{R}$, $a[i] \in \mathbb{R}$, $a(i) \in \mathbb{R}$ Several styles
 - ✓ Natural choice: a_i
 - ✓ Don't write: \mathbf{a}_i (this implies an i -th vector from $\{\mathbf{a}_n\}_{n=1}^{10}$)

- Accessing an element of a matrix: Given $\mathbf{A} \in \mathbb{R}^{3 \times 2}$,
 - ✓ (i, j) -th element: $A_{ij} \in \mathbb{R}$, $a_{ij} \in \mathbb{R}$, $A[i, j] \in \mathbb{R}$, $A(i, j) \in \mathbb{R}$ Several styles
 - ✓ i -th row: $\mathbf{A}_{i:} \in \mathbb{R}^{1 \times 2}$, $\mathbf{A}[i, :] \in \mathbb{R}^{1 \times 2}$, $\mathbf{A}(i, :) \in \mathbb{R}^{1 \times 2}$
 - ✓ j -th column: $\mathbf{A}_{:, j} \in \mathbb{R}^3$, $\mathbf{A}[:, j] \in \mathbb{R}^3$, $\mathbf{A}(:, j) \in \mathbb{R}^3$ One may not use bold

Matrix and tensor

- Parenthesis or square brackets?

- ✓ Both are fine.

$$A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix} \quad B = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$$

- Personally, I prefer square brackets as parenthesis have more meanings.

- i.e., parenthesis for vector/matrix may conflict with other usages, e.g.,

- ✓ Usual sentence: "We think XXX (Note that $\mathbf{x} = (1, 2)$)"

- ✓ Binomial coefficient: $\binom{2}{3}$ Cannot distinguish with $\binom{2}{3} \in \mathbb{R}^2$

- ✓ Element-based representation of a matrix: $\mathbf{A} = (a_{ij})$

- ✓ Function: $f((1, 2))$

- ✓ Tuple: $(1, 2, 3)$ Parenthesis is used for a tuple.

- ✓ etc...

Set

- Recommend: Calligraphy font in uppercase, e.g.,
 - ✓ $\mathcal{S} = \{2, 13, 5, 7\} \subset \mathbb{N}$ A set of natural numbers.
 - ✓ $\mathcal{X} = \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N\} \subset \mathbb{R}^D$ A set of N D -dimensional vectors.
 - Can be written as $\mathcal{X} = \{\mathbf{x}_n\}_{n=1}^N$. Here, $\mathbf{x}_n \in \mathbb{R}^D$.
- How to write
 - ✓ TeX: `\mathcal{A}`
 - ✓ Powerpoint: e.g., `\scriptS`
- Cardinality (the number of elements): $|\mathcal{S}| = 4$
- A set doesn't contain duplicate elements.
 - ✓ If does, it's called multi-set (or bag): $\mathcal{A} = \{a, a, b\}$.

Set

➤ Domain??

- ✓ 😞 It's a bit tough to show a domain of a set.

Powerset: Given a set \mathcal{A} , a powerset of \mathcal{A} is defined as all subsets of \mathcal{A} . E.g.,

➤ $\mathcal{A} = \{1, 3, 5\}$, then

➤ $2^{\mathcal{A}} = \{\phi, \{1\}, \{3\}, \{5\}, \{1, 3\}, \{3, 5\}, \{5, 1\}, \{1, 3, 5\}\}$

- ✓ Consider a set $\mathcal{S} = \{5, 3\} \subset \mathcal{A}$, the domain of a set \mathcal{S} is:

□ $\mathcal{S} \in 2^{\mathcal{A}}$

- ✓ In the same way, if a set \mathcal{B} is a subset of \mathcal{X} , \mathcal{B} 's domain is $2^{\mathcal{X}}$

□ $\mathcal{B} \subset \mathbb{R}$, then $\mathcal{B} \in 2^{\mathbb{R}}$

Usually, " $\mathcal{B} \subset \mathbb{R}$ " style is easier to read.

➤ The powerset may helps when defining a function with a set-input $f(\mathcal{B})$:

✓ $f: 2^{\mathbb{R}} \rightarrow \mathbb{R}$

Summary (again)

| Type | How to write | Example of domain | Example of values |
|-------------------|--|--|--|
| Scalar | Lowercase or uppercase | $a \in \mathbb{R}. b \in \mathbb{N}. K \in \{10, 20, 30\}.$ | $a = 3.2. b = 13. K = 20.$ |
| Vector | Lowercase and bold ➤ TeX: <code>\mathbf</code> or <code>\bm</code> ➤ Powerpoint: bold | $\mathbf{x} \in \mathbb{R}^3. \mathbf{b} \in \{0, 1\}^B.$ | $\mathbf{x} = [0.1, 0.2, 0.3]^\top. \mathbf{b} = [0, 1, 1, 0]^\top.$ |
| Matrix and Tensor | Uppercase and bold ➤ TeX: <code>\mathbf</code> or <code>\bm</code> ➤ Powerpoint: bold | $\mathbf{A} \in \mathbb{R}^{2 \times 3}. \mathbf{I} \in [0, 1]^{H \times W \times 3}.$ | $\mathbf{A} = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 4 \end{bmatrix}.$ |
| Set | Calligraphy font in uppercase ➤ TeX: <code>\mathcal</code> ➤ Powerpoint: e.g., <code>\scriptS</code> | $\mathcal{S} \subset \mathbb{N}.$ | $\mathcal{S} = \{2, 4, 8, 16\}.$ |

Equations

- Basic notation for variables
- **String**
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

String 🤔 It's not easy to use strings in equations

- If you use strings for labels, the easiest way is to define them directly.
 - ✓ A label set $\mathcal{L} = \{\text{"dog"}, \text{"cat"}, \text{"horse"}\}$
 - ✓ Image classifier $f: \mathbb{R}^{H \times W \times 3} \rightarrow \mathcal{L}$
- Alternatively, you can assign an integer for each label.
 - ✓ A label set $\mathcal{L} = \{1, \dots, N\}$
 - ✓ Image classifier $f: \mathbb{R}^{H \times W \times 3} \rightarrow \mathcal{L}$
- Alternatively, you can use one-hot vectors.
 - ✓ A label set $\mathcal{L} = \{\mathbf{b}_n\}_{n=1}^N, \mathbf{b}_n \in \{0, 1\}^N$
 - ✓ Image classifier $f: \mathbb{R}^{H \times W \times 3} \rightarrow \mathcal{L}$
 - ✓ Easy to integrate in other equations, e.g., $\|f(\mathbf{X}) - \mathbf{y}\|_2^2$ for some \mathbf{y}

If you don't run more complex operations on them, this should suffice.

| Label | ID | One-hot |
|---------|----|---------------|
| "dog" | 1 | $[1, 0, 0]^T$ |
| "cat" | 2 | $[0, 1, 0]^T$ |
| "horse" | 3 | $[0, 0, 1]^T$ |

- The description $\mathbf{b}_n \in \{0, 1\}^N$ is not tight; could be any N -dim binary vectors.
- If you want to emphasize \mathbf{b}_n is one-hot, you can write:
 - ✓ $\mathbf{b}_n \in \mathcal{H}_{1/N}$
 - ✓ $\mathcal{H}_{1/N} = \{\mathbf{b} \in \{0, 1\}^N \mid \|\mathbf{b}\| = 1\}$
 - ✓ Here, $\mathcal{H}_{1/N}$ is a set of all 1-of- N binary encoding vectors [1]

[1] M. Norouzi and D. J. Fleet, "Cartesian k-means", CVPR 2013

| Label | ID | One-hot |
|---------|----|---------------|
| "dog" | 1 | $[1, 0, 0]^T$ |
| "cat" | 2 | $[0, 1, 0]^T$ |
| "horse" | 3 | $[0, 0, 1]^T$ |

- Alternatively, you can use one-hot vectors
 - ✓ A label set $\mathcal{L} = \{\mathbf{b}_n\}_{n=1}^N, \mathbf{b}_n \in \{0, 1\}^N$
 - ✓ Image classifier $f: \mathbb{R}^{H \times W \times 3} \rightarrow \mathcal{L}$
 - ✓ Easy to integrate in other equations, e.g., $\|f(\mathbf{X}) - \mathbf{y}\|_2^2$ for some \mathbf{y}

String AB A set of character sequences \mathcal{V}^*

- If you need more detailed definition: Kleene closure for characters.
 - ✓ Define a vocabulary $\mathcal{V} = \{\text{“a”}, \text{“b”}, \dots, \text{“z”}\}$.
 - ✓ Here, $\mathcal{V}^2 = \mathcal{V} \times \mathcal{V}$, $\mathcal{V}^3 = \mathcal{V} \times \mathcal{V} \times \mathcal{V}$, ... i.e., $\text{“a”} \in \mathcal{V}$ and $\text{“dog”} \in \mathcal{V}^3$
 - ✓ Define **a set of character sequences** $\mathcal{V}^* = \{\emptyset\} \cup \mathcal{V} \cup \mathcal{V}^2 \cup \dots$
 - ✓ Any string (character sequence) is in \mathcal{V}^* : $\text{“a”}, \text{“dog”}, \text{“hoge”} \in \mathcal{V}^*$

- A label set $\mathcal{L} = \{\text{“dog”}, \text{“cat”}, \text{“horse”}\} \subset \mathcal{V}^*$

- Straightforward. But not easy to connect with other equations.

String 1 2 3 4 A set of token sequences \mathcal{V}^*

| char | token |
|------|-------|
| "a" | 1 |
| "b" | 2 |
| ... | |
| "z" | N |

N : vocabulary size

➤ If you need more math-friendly representations:
Kleene closure for tokens.

✓ Define a vocabulary $\mathcal{V} = \{1, \dots, N\}$.

✓ Here, $\mathcal{V}^2 = \mathcal{V} \times \mathcal{V}$, $\mathcal{V}^3 = \mathcal{V} \times \mathcal{V} \times \mathcal{V}$, ...

✓ i.e., $3 \in \mathcal{V}$ and $[13, 6, 20] \in \mathcal{V}^3$

➤ Now it's just a vector.

➤ Intentionally use "row-vector" style

✓ Define **a set of token sequences** $\mathcal{V}^* = \{\emptyset\} \cup \mathcal{V} \cup \mathcal{V}^2 \cup \dots$

✓ Any token sequence (variable length integer-vectors) is in \mathcal{V}^*

✓ $3, [13, 6, 20], [6, 18, 16, 7] \in \mathcal{V}^*$

➤ This is an ASCII representation for the c-language.

➤ [2] uses this notation.

String 1 2 3 4 A set of token sequences

| v | $D(v)$ |
|-----|--------|
| 1 | “a” |
| 2 | “b” |
| ... | ... |
| N | “z” |

- Decoder
 - ✓ Given $v \in \mathcal{V}$, decoder $D(v)$ returns the original character.
 - ✓ e.g., $D(7) = \text{“g”}$.
 - ✓ Extends this to $\mathbf{w} \in \mathcal{V}^*$. e.g., $D([4, 15, 7]) = \text{“dog”}$.

- Vector-style operations
 - ✓ e.g., $\mathbf{w} = [6, 18, 16, 7] \in \mathcal{V}^*$, then $D(\mathbf{w}) = \text{“frog”}$
 - ✓ Specify a character: $w_2 = 18$, and $D(w_2) = \text{“r”}$
 - ✓ Slicing: $\mathbf{w}_{2:4} = [18, 16, 7]$, and $D(\mathbf{w}_{2:4}) = \text{“rog”}$

- Difference to a superset? $2^{\mathcal{V}}$ vs \mathcal{V}^*
 - ✓ An element of a super set is a set (not a vector). No duplicate.
 - ✓ $\{3, 4\} \in 2^{\mathcal{V}}$ Cannot have duplicates $[3, 4, 4] \in \mathcal{V}^*$ Can have duplicates

String 1 2 3 4 A set of token sequences

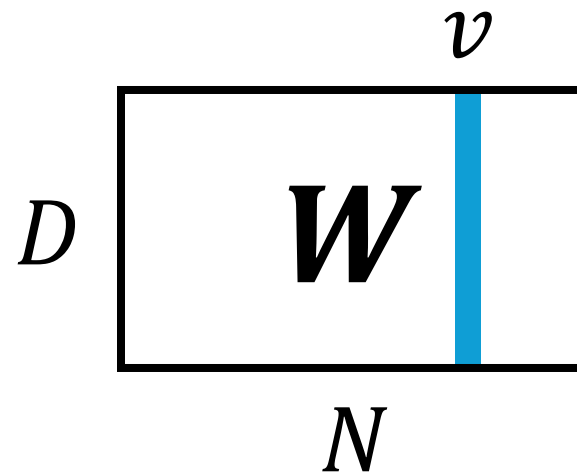
➤ Token embedding

✓ Embedding matrix $\mathbf{W} \in \mathbb{R}^{D \times N}$

D : dim of embedding

✓ The vector representation (embedding) of $v \in \mathcal{V}$ is $\mathbf{W}[:, v] \in \mathbb{R}^D$

| v | $D(v)$ |
|-----|--------|
| 1 | "a" |
| 2 | "b" |
| ... | ... |
| N | "z" |



➤ Can be a tuple?

✓ Yes. You can define so (I'll discuss later)

✓ $(1, 3, 5) \in \mathcal{V}^*$

Equations

- Basic notation for variables
- String
- **Tips for sets**
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

Tips for sets

➤ How to describe $\{x_1, x_2, \dots, x_N\}$ in short?

✓ 😐 OK, but a bit long? $\{x_i \mid i = 1, 2, \dots, N\}$

✓ 😐 OK, but a bit long? $\{x_i \mid 1 \leq i \leq N\}$

✓ 😐 OK, but a bit long? $\{x_i \mid i \in \{1, 2, \dots, N\}\}$

✓ 😊 Good! Short! $\{x_i\}_{i=1}^N$

This may seem obvious? But many junior students struggle with this.

➤ How to do when I don't want to use an alphabet for #set?

✓ $\mathcal{X} = \{x_1, x_2, \dots\}$ — Tips to hide the number of elements

✓ If you need the number of \mathcal{X} , use $|\mathcal{X}|$

You don't consume any additional alphabets

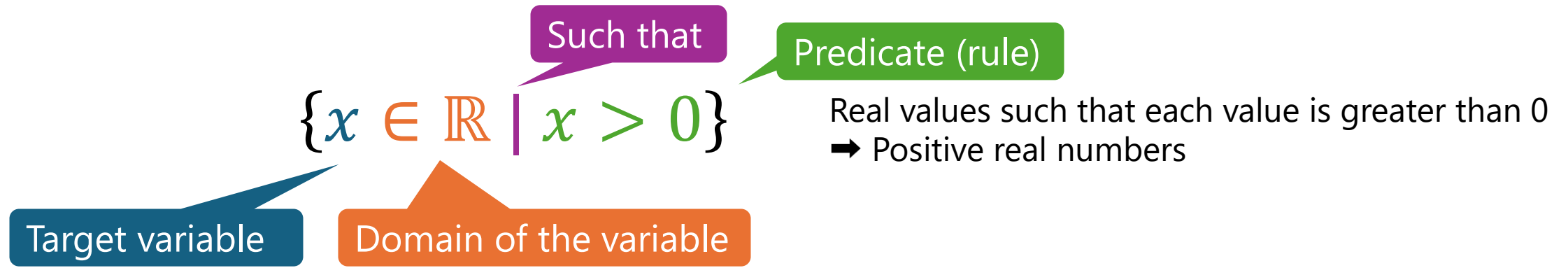
✓ Double brackets to enumerate natural numbers

✓ $\llbracket N \rrbracket = \{1, 2, \dots, N\}$

✓ 🤔 This form is simple, but not so much intuitive. Be careful.

Set-builder notation

- A way to create a set. e.g.,



- The domain can be written on the right:

✓ $\{x \in \mathbb{R} \mid x > 0\} = \{x \mid x \in \mathbb{R}, x > 0\}$

- Any function can be applied to the target variable:

✓ $\{3x + 5 \mid x \in \mathbb{R}^4, \|x\|_2 = 1\}$ For all points in the 4-dimensional unit sphere, affine-transform them by scaling 3 and shift 5

Set-builder notation

- Multiple predicates
 - ✓ $\{\mathbf{x} \in \mathbb{R}^3 \mid \mathbf{x}^\top \mathbf{y}_1 = 0, \mathbf{x}^\top \mathbf{y}_2 = 0\}$ All 3-dim vectors that are orthogonal to both \mathbf{y}_1 and \mathbf{y}_2
- Multiple variables can be used at the same time (multiple for-loop)
 - ✓ $\{5ij \mid i, j \in \{1, 2, 3\}, i \neq j\}$
- Set-builder notation is equivalent to a list comprehension in Python
 - ✓ $\mathcal{A} = \{1, 2, 3, 4, 5\}$, and $\{a^2 \mid a \in \mathcal{A}, a > 3\}$
 - ✓ $A = [1, 2, 3, 4, 5]$ and $\{a*a \text{ for } a \text{ in } A \text{ if } a > 3\}$

Equations

- Basic notation for variables
- String
- Tips for sets
- **Inputs/outputs of functions**
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

Inputs/outputs of functions

- There are two ways to describe the inputs/outputs of a function
 - ✓ e.g., considering $f(x) = x^2$ for $x \in \mathbb{R}$

$f: \mathbb{R} \rightarrow \mathbb{R}$ Set-based description. Use `\to` (\rightarrow)

$f: x \mapsto x^2$ Element-based description. Use `\mapsto` (\mapsto)

- Describing the inputs/outputs makes functions easier to read.
- It's ok to directly write the description in a sentence. **No side effects.**
 - ✓ 😐 "Let us define a function f as follows ..."
 - ✓ 😊 "Let us define a function $f: \mathbb{N} \rightarrow \{1, \dots, 10\}$ as follows ..."

Inputs/outputs of functions

Domain

Codomain

There are two ways to describe the inputs/outputs of a function

✓ e.g. considering $f(x) = x^2$ for $x \in \mathbb{R}$

$$f: \mathbb{R} \rightarrow \mathbb{R}$$
$$f: x \mapsto x^2$$

Set-based description. Use `\to` (\rightarrow)

Element-based description. Use `\mapsto` (\mapsto)

- Describing the inputs/outputs makes functions easier to read.
- It's ok to directly write the description in a sentence. **No side effects.**
 - ✓ 😐 "Let us define a function f as follows ..."
 - ✓ 😊 "Let us define a function $f: \mathbb{N} \rightarrow \{1, \dots, 10\}$ as follows ..."

Inputs/outputs of functions

- Multiple inputs: $z = f(x, y) = x^2 + y + 1$
 - ✓ $f: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$
 - ✓ $f: (x, y) \mapsto z$ or $f: (x, y) \mapsto x^2 + y + 1$
- Vector input: $f(\mathbf{x}) = \mathbf{a}^\top \mathbf{x} + b$, where $\mathbf{a}, \mathbf{x} \in \mathbb{R}^D$ and $b \in \mathbb{R}$
 - ✓ $f: \mathbb{R}^D \rightarrow \mathbb{R}$
 - ✓ $f: \mathbf{x} \mapsto \mathbf{a}^\top \mathbf{x} + b$
- Vector input and vector output: for $\mathbf{x} = [x_1, x_2]^\top \in \mathbb{R}^2$, f is defined as
 - ✓ $f: \mathbf{x} \mapsto \begin{bmatrix} x_1 + x_2 \\ 3x_1 + \log x_2 \\ x_2^3 \end{bmatrix}$ by the element-based description.
 - ✓ Then, the set-based description is $f: \mathbb{R}^2 \rightarrow \mathbb{R}^3$

Inputs/outputs of functions

- When the output is a vector, should the function itself be bold?
 - ✓ Depends.

$$f: \mathbb{R}^2 \rightarrow \mathbb{R}^3$$

$$\mathbf{f}: \mathbb{R}^2 \rightarrow \mathbb{R}^3$$

- 😊 With this, the equation may be beautiful: $\mathbf{z} = \mathbf{f}(\mathbf{x}) + \mathbf{a}$
- 🤔 In modern CV, all functions may have a vector-output, thus all functions may be bold. It may seem "heavy"?
- In this document, I use non-bold for vector-functions.

Set-based? Element-based?

- Usually, the set-based description is more informative.
- Readers have an interest in what are the possible values of the inputs and outputs.

- e.g., for $f(x) = x^2$,

$$f: \mathbb{R} \rightarrow \mathbb{R}$$

$$f: x \mapsto x^2$$

😊 Both an input and an output are real-values.

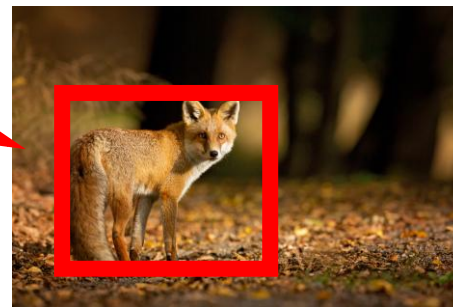
😐 This is just repeating the definition.

- The element-based description is useful if you want to **intentionally hide** the complex relationships.

Set-based? Element-based?

- Consider an object detector f , which inputs $H \times W \times 3$ 8-bit image (each pixels is in $\{0, \dots, 255\}$) and returns the followings.

Label $l \in \{1, \dots, K\}$
Bounding box $\mathbf{b} = [y, x, h, w]^T \in \mathbb{N}^4$
Confidence score $\alpha \in \mathbb{R}$



- Set-based precise description is complex:
 - ✓ $f: \{0, \dots, 255\}^{H \times W \times 3} \rightarrow \{1, \dots, K\} \times \mathbb{N}^4 \times \mathbb{R}$

😬 The pixel value range is not important.
Don't want to use alphabets H and W here.

😬 Not so much clear about the right side.
What is each variable?

- May be better to moderately hide the detail:

"We consider a K -class object detector $f: I \mapsto (l, \mathbf{b}, \alpha)$. This function inputs an image I , and returns a label $l \in \{1, \dots, K\}$, a bbox $\mathbf{b} = [y, x, h, w]^T \in \mathbb{N}^4$, and a confidence $\alpha \in \mathbb{R}$ "

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- **Subscript/superscript**
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- Misc

Subscript/superscript

- Superscripts and subscripts are used to provide additional information to a variable.



- Principal: **As little use as possible!**
- $x_{i,j}^k$ Three variables: 😞 No! Hard to follow...
- x_{a_i} Subscript of subscript: 😞 No! Hard to follow...
- Alternatives for subscript/superscript: decorations
 - e.g., \hat{x} , \bar{x}
 - No: $x_{\text{mean}} = \frac{1}{N} \sum x_n$ Recommend: $\bar{x} = \frac{1}{N} \sum x_n$

Subscript/superscript

- Tips: a loop index can be removed by re-define a variable.
- Consider $\mathcal{V} = \{v_n\}_{n=1}^N$

- 😐 "Here, $f(v_n)$ is ..."

n is not important

```
for n in range(len(V)):  
    f(V[n])
```

- 😊 "Let us consider $v \in \mathcal{V}$. Here, $f(v)$ is ..."

Remove it!

```
for v in V:  
    f(v)
```

Subscript/superscript

- Don't make the subscript bold (very typical mistake)

✓ 😞 `\mathbf{x}_i`: x_i

✓ 😊 `\mathbf{x}_i`: x_i

- Natural interpretation is that i is a vector to indicate identifiers, e.g., $i = [1, 2]$, results in $x_{1,2}$
- Moreover, this mistake smells amateurish.

- If two+ letters are used in the meaning of a label, make it Roman.

✓ `\mathbf{x}_{\mathrm{in}}`: x_{in} OK


✓ `\mathbf{x}_{in}`: x_{in} Not recommend

- This can be interpreted as

✓ $j = i * n$

✓ x_j

Subscript/superscript

- Don't put variables with different meanings in sub/superscript.
- Consider $\mathcal{V} = \{\mathbf{v}_n\}_{n=1}^N$, where $\mathbf{v}_n \in \mathbb{R}^D$
- How to denote d -th element of n -th vector?
 - ✓ 😞 $v_{n,d} \in \mathbb{R}$ Not recommend! n and d have different meanings!
 - ✓ 😬 $v_n[d] \in \mathbb{R}$: OK. Combining square brackets.
 - ✓ 😊 $v_d \in \mathbb{R}$ or $v[d] \in \mathbb{R}$ for $\mathbf{v} \in \mathcal{V}$: Clear.
- Tips: You can stack \mathcal{V} to form a matrix $\mathbf{V} \in \mathbb{R}^{D \times N}$ 
 - ✓ Then, $v_{d,n} \in \mathbb{R}$ or $v[d,n] \in \mathbb{R}$ is OK (usual matrix element access)

(n, d) is now (d, n). Be careful!
- Vector-set (\mathcal{V}) to matrix (\mathbf{V}) doesn't consume an alphabet.

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- **Don't write numpy**
- Misc

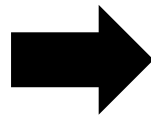
Pseudo codes

- Basic
- What do you want to express?
- Misc

Don't write numpy

- Many people write numpy descriptions directly in equations. It's wrong.

```
a = np.array([5, 4, 3])  
b = a - 3
```



```
a = [5, 4, 3]  
b = a - 3
```

😞 A horrible mistake!

- Numpy can broadcast a scalar
- We can get $b = [2, 1, 0]$

- Math equations don't broadcast.
- This operation is undefined!

- Recommend:

- $\mathbf{a} = [5, 4, 3]^T$. $\mathbf{b} = \mathbf{a} - 3\mathbf{1}$

- Use bold
- Column-vectors

- Use all-one vector $\mathbf{1} = [1, 1, \dots]^T$

Specific examples of failures

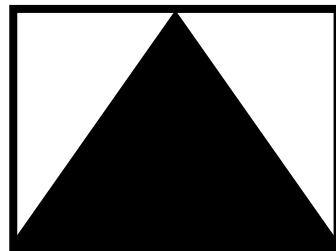
- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = BX + (1 - B)Y$$



Z

=

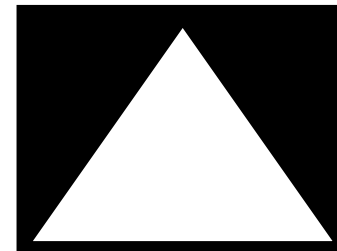


B



X

+



$1 - B$



Y

There are three big mistakes... Can you guess? 🤔

Specific examples of failures

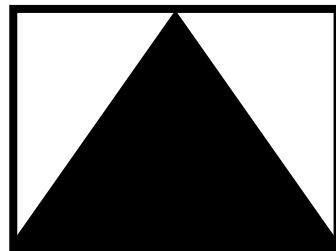
- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$\mathbf{Z} = \mathbf{B}\mathbf{X} + (1 - \mathbf{B})\mathbf{Y}$$



$$\mathbf{Z} \in \mathbb{R}^{H \times W \times 3}$$

=

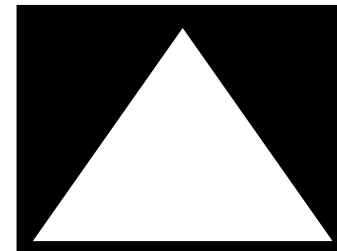


$$\mathbf{B} \in \{0, 1\}^{H \times W}$$



$$\mathbf{X} \in \mathbb{R}^{H \times W \times 3}$$

+



$$1 - \mathbf{B}$$



$$\mathbf{Y} \in \mathbb{R}^{H \times W \times 3}$$

First, make the variables bold, and write domains

Specific examples of failures

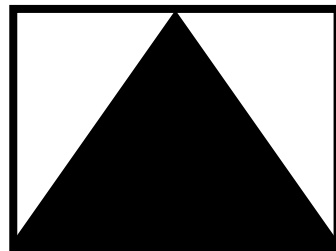
- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$\mathbf{Z} = \mathbf{B}\mathbf{X} + (1 - \mathbf{B})\mathbf{Y}$$



$$\mathbf{Z} \in \mathbb{R}^{H \times W \times 3}$$

=

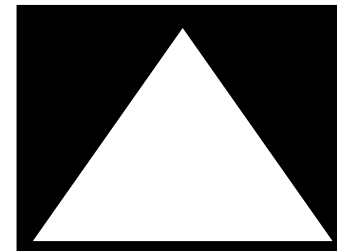


$$\mathbf{B} \in \{0, 1\}^{H \times W}$$



$$\mathbf{X} \in \mathbb{R}^{H \times W \times 3}$$

+



$$1 - \mathbf{B}$$



$$\mathbf{Y} \in \mathbb{R}^{H \times W \times 3}$$

First, make the variables bold, and write domains

Mistake 1: Broadcast happens here!
We need to use an all-one matrix.

Specific examples of failures

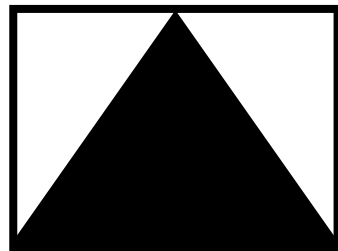
- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = BX + (1 - B)Y$$



$$Z \in \mathbb{R}^{H \times W \times 3}$$

=

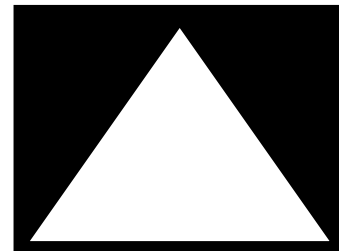


$$B \in \{0, 1\}^{H \times W}$$



$$X \in \mathbb{R}^{H \times W \times 3}$$

+



$$1 - B$$



$$Y \in \mathbb{R}^{H \times W \times 3}$$

Specific examples of failures

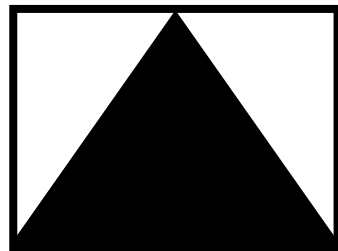
- Let X and Y be 3-channel images.
- **Mistake2: BX is matrix-multiplication! We need per-element multiplication (Hadamard product) \odot** when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = BX + (1 - B)Y$$



$$Z \in \mathbb{R}^{H \times W \times 3}$$

=

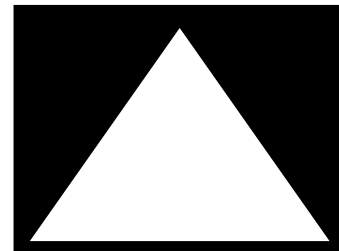


$$B \in \{0, 1\}^{H \times W}$$



$$X \in \mathbb{R}^{H \times W \times 3}$$

+



$$1 - B$$



$$Y \in \mathbb{R}^{H \times W \times 3}$$

Specific examples of failures

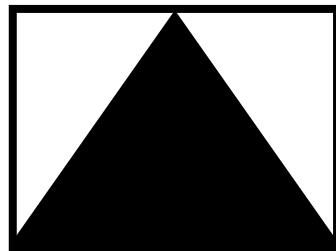
- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



$$Z \in \mathbb{R}^{H \times W \times 3}$$

=



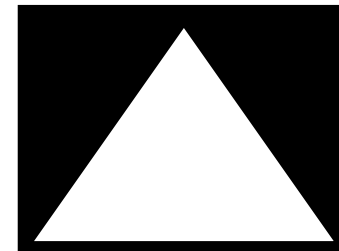
$$B \in \{0, 1\}^{H \times W}$$

\odot



$$X \in \mathbb{R}^{H \times W \times 3}$$

+



$$1 - B$$

\odot




$$Y \in \mathbb{R}^{H \times W \times 3}$$

Specific examples of failures

- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



$Z \in \mathbb{R}^{H \times W \times 3}$ $B \in \{0, 1\}^{H \times W}$ $X \in \mathbb{R}^{H \times W \times 3}$ $1 - B$ $Y \in \mathbb{R}^{H \times W \times 3}$

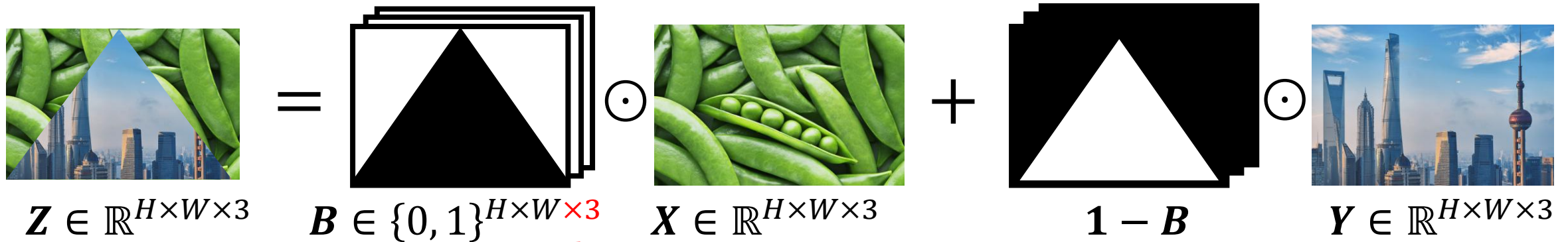
Mistake3: $B \odot X$ still doesn't work because the tensor shape is different! $H \times W$ vs $H \times W \times 3$

👤 80% computer vision papers ignore this issue...

Specific examples of failures

- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



Solution 1: Define B as a stack of the original B , making the shape same.

Specific examples of failures

- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



The diagram illustrates the failure of element-wise operations on images. It shows the equation $Z = B \odot X + (1 - B) \odot Y$. The resulting image Z is a combination of X and Y , but the mask B is not a binary image, leading to a failure in the element-wise operations. The diagram shows the resulting image Z (a cityscape with a red 'X' over the dimension notation), the mask B (a black triangle on a white background), the image X (peas), the mask $1 - B$ (a white triangle on a black background), and the image Y (a cityscape with a red 'X' over the dimension notation). The dimension notation for Z and Y is $\mathbb{R}^{H \times W \times 3}$.

Solution 2: Describe everything per channel, then combine the results.

Specific examples of failures

- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



Solution 3: Define \odot loosely, e.g.,

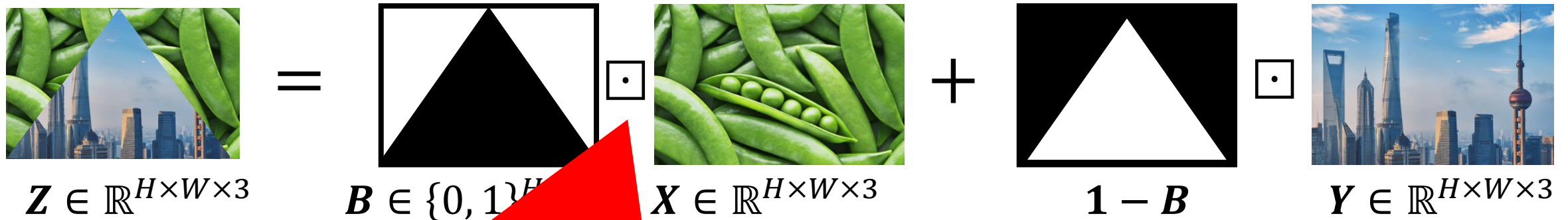
²Similar to many papers, we also abuse this operator slightly; $X \odot (1 - t(\phi(V)))$ computes the element-wise product between $(1 - t(\phi(V)))$ and each channel of X .

<https://arxiv.org/abs/2311.15994>

Specific examples of failures

- Let X and Y be 3-channel images.
- Consider a mask matrix B whose elements are 0 or 1.
- We adopt X when the value of the mask is one and Y when the value of the mask is zero.
- Z , the resulting image combining X and Y , is computed as follows.

$$Z = B \odot X + (1 - B) \odot Y$$



Solution 4: Use broadcast product \odot

Y. Matsui & T. Yokota, "Broadcast Product: Shape-aligned Element-wise Multiplication and Beyond", arXiv 2024
<https://arxiv.org/abs/2409.17502>

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- **Misc**

Pseudo codes

- Basic
- What do you want to express?
- Misc

Misc

- Don't write English words in equations

- 😞 $y = score(x) + 10$

- 😊 $y = s(x) + 10$

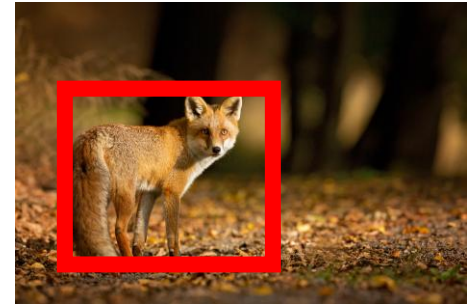
➤ This can be interpreted as $s * core(x)$

- If you're not 100% confident, don't use quantifiers (\forall, \exists).
- If you use variables, please always define and explain them. E.g., if you write $y = ax + b$, then explain y, a, x, b .
- If it is too difficult to describe what you are trying to explain in mathematical equations, please explain them in writing and figures.
 - ✓ 😞 Writing half-wrong equations are terrible.

Tuple 🍡

- (something like) an ordered-set: e.g., $a = (10, 20, 30)$
- Similar to a set:
 - ✓ Can contain any elements: e.g., $b = ([1, 2, 3]^T, "a", 25)$
 - ✓ Remember: $f: I \mapsto (l, \mathbf{b}, \alpha)$
- But the order is decided:
 - ✓ $(1, 2, 3) \neq (2, 1, 3)$
 - ✓ $\{1, 2, 3\} = \{2, 1, 3\}$
- Can contain duplicates: $a = (1, 1, 3)$
- Similar to a vector:
 - ✓ $\mathbf{c} = [1, 2, 3]^T$ vs $d = (1, 2, 3)$
- But (usually) no mathematical structures:
 - ✓ E.g., the addition is not defined: $(1, 2, 3) + (4, 5, 6)$

This was a tuple



No!

Cheat-sheet

| Not-recommend | OK | Reason |
|---|--|---|
| Write a vector as x | Write a vector as \boldsymbol{x} | Use a bold font for a vector. |
| $\boldsymbol{x} = [1, 2, 3]$ | $\boldsymbol{x} = [1, 2, 3]^T$ | Use column-vectors. |
| Consider D -dim vector x | Consider D -dim vector $\boldsymbol{x} \in \mathbb{R}^D$ | Show the domain. |
| $x_{i,j}^k$ Many sub-/superscripts | Don't use so much. | Hard to understand. |
| x_{a_i} subscript of subscript | Avoid. | Hard to understand. |
| \boldsymbol{x}_i | \boldsymbol{x}_i | Don't make an index bold. |
| x_{in} | x_{in} | Labels should be roman. |
| $score(x) + 10$ | $s(x) + 10$ | Don't use English words. |
| $a \in \mathbb{R}^3, b \in \mathbb{R}$, then $a + b$ | $\boldsymbol{a} + b\mathbf{1}$ | Don't broadcast. |
| For element-wise product, \boldsymbol{AB} | $\boldsymbol{A} \odot \boldsymbol{B}$ | Element-wise product is Hadamard product. |
| Using \forall, \exists but not 100% confident | Don't use | You are wrong. |

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- **Basic**
- What do you want to express?
- Misc

Pseudo-code: Basic

- What is pseudo-code?
 - ✓ Describe an algorithm
 - ✓ OK to use mathematical equations
 - ✓ OK to use data structures
 - ✓ Highlight the key idea
 - ✓ Many ways to explain the idea

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$  for  $n = 1, \dots, N$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
    
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

Algorithm 3 SVD1 : QB-backed low-rank SVD (see [HMT11] and [RST10])

```

1: function SVD1( $\mathbf{A}, k, \epsilon, s$ )
  Inputs:
     $\mathbf{A}$  is an  $m \times n$  matrix. The returned approximation will have rank at most  $k$ . The approximation produced by the randomized phase of the algorithm will attempt to  $\mathbf{A}$  to within  $\epsilon$  error, but will not produce an approximation of rank greater than  $k + s$ .
  Output:
    The compact SVD of a low-rank approximation of  $\mathbf{A}$ .
  Abstract subroutines:
    QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.
2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon) \# \mathbf{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, : r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, : r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, r, : r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
    
```

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- What is pseudo-code?
 - ✓ Describe an algorithm
 - ✓ OK to use mathematical equations
 - ✓ OK to use data structures
 - ✓ Highlight the key idea
 - ✓ Many ways to explain the idea

Inputs / outputs description

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$  for  $n = 1, \dots, N$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
    
```

Algorithm 3 SVD-based low-rank SVD (see [HMT11] and [RST10])

```

1: function SVD( $\mathbf{A}, k, \epsilon, s$ )
    
```

Inputs:
 \mathbf{A} is an $m \times n$ matrix. The returned approximation will have rank at most k . The approximation produced by the randomized phase of the algorithm will attempt to approximate \mathbf{A} to within ϵ error, but will not produce an approximation of rank greater than $k + s$.

Output:
 The compact SVD of a low-rank approximation of \mathbf{A} .

Abstract subroutines:
 QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.

```

2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon) \# \mathbf{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, : r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, : r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, r, : r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
    
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- What is pseudo-code?
 - ✓ Describe an algorithm
 - ✓ OK to use mathematical equations
 - ✓ OK to use data structures
 - ✓ Highlight the key idea
 - ✓ Many ways to explain the idea

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
    
```

Can run other functions or pseudo-codes

Algorithm 3 SVD1 : QB-backed low-rank SVD (see [HMT11] and [RST10])

1: function SVD1($\mathbf{A}, k, \epsilon, s$)

Inputs:

\mathbf{A} is an $m \times n$ matrix. The returned approximation will have rank at most k . The approximation produced by the randomized phase of the algorithm will attempt to approximate \mathbf{A} to within ϵ error, but will not produce an approximation of rank greater than $k + s$.

Output:

The compact SVD of a low-rank approximation of \mathbf{A} .

Abstract subroutines:

QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.

```

2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon)$  #  $\mathbf{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, : r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, : r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, : r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
    
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- What is pseudo-code?
 - ✓ Describe an algorithm
 - ✓ OK to use mathematical equations
 - ✓ OK to use data structures
 - ✓ Highlight the key idea
 - ✓ Many ways to explain the idea

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$  for  $n = 1, \dots, N$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
    
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

Algorithm 3 SVD1 : QB-backed low-rank SVD (see [HMT11] and [RST10])

1: **function** SVD1($\mathbf{A}, k, \epsilon, s$)

Inputs:

\mathbf{A} is an $m \times n$ matrix. The returned approximation will have rank at most k . The approximation produced by the randomized phase of the algorithm will attempt to approximate \mathbf{A} to within ϵ error, but will not produce an approximation of rank greater than $k + s$.

Output:

The compact SVD of a low-rank matrix.

Abstract subroutines:

QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.

```

2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon)$  #  $\text{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, :r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, :r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, r, :r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
    
```

Comments

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- What is pseudo-code?
 - ✓ Describe an algorithm
 - ✓ OK to use mathematical equations
 - ✓ OK to use data structures

- Can use usual sentences to represent complex operations.
- If you think it's hard to write your operations by equations, use sentences.

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$  for  $n = 1, \dots, N$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
    
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

Algorithm 3 SVD1 : QB-backed low-rank SVD (see [HMT11] and [RST10])

1: **function** SVD1($\mathbf{A}, k, \epsilon, s$)

Inputs:

The compact SVD of a low-rank approximation of \mathbf{A} .

Abstract subroutines:

QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.

```

2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon) \# \mathbf{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, : r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, : r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, r, : r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
    
```

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- Look very different
- Left: more pseudo-code-ish
- Right: more like equations (this is actually MATLAB-ish)

- ✓ OK to use mathematical structures
- ✓ Highlight the main idea
- ✓ Many ways to explain the idea

```

procedure CP-ALS( $\mathcal{X}, R$ )
  initialize  $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times R}$  for  $n = 1, \dots, N$ 
  repeat
    for  $n = 1, \dots, N$  do
       $\mathbf{V} \leftarrow \mathbf{A}^{(1)\top} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)\top} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)\top} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)\top} \mathbf{A}^{(N)}$ 
       $\mathbf{A}^{(n)} \leftarrow \mathbf{X}^{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)}) \mathbf{V}^\dagger$ 
      normalize columns of  $\mathbf{A}^{(n)}$  (storing norms as  $\lambda$ )
    end for
  until fit ceases to improve or maximum iterations exhausted
  return  $\lambda, \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)}$ 
end procedure
  
```

Fig. 3.3 ALS algorithm to compute a CP decomposition with R components for an N th-order tensor \mathcal{X} of size $I_1 \times I_2 \times \dots \times I_N$.

Kolda and Bader, "Tensor Decompositions and Applications", SIAM Review, 2009.

SVD (see [HMT11] and [RST10])

returned approximation will have rank at most $k + s$. The approximation produced by the randomized phase of the algorithm is ϵ -close to \mathbf{A} to within ϵ error, but will not produce an approximation of rank greater than $k + s$.

Output: $\mathbf{Q}, \mathbf{\Sigma}, \mathbf{V}^*$
 The decomposition of a low-rank approximation of \mathbf{A} .

Abstract syntax:

QBDecomposer generates a QB decomposition of a given matrix; it tries to reach a prescribed error tolerance but may stop early if it reaches a prescribed rank limit.

```

2:  $\mathbf{Q}, \mathbf{B} = \text{QBDecomposer}(\mathbf{A}, k + s, \epsilon) \# \mathbf{QB} \approx \mathbf{A}$ 
3:  $r = \min\{k, \text{number of columns in } \mathbf{Q}\}$ 
4:  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^* = \text{svd}(\mathbf{B})$ 
5:  $\mathbf{U} = \mathbf{U}[:, :r]$ 
6:  $\mathbf{V} = \mathbf{V}[:, :r]$ 
7:  $\mathbf{\Sigma} = \mathbf{\Sigma}[:, :r, :r]$ 
8:  $\mathbf{U} = \mathbf{Q}\mathbf{U}$ 
9: return  $\mathbf{U}, \mathbf{\Sigma}, \mathbf{V}^*$ 
  
```

Murray+, "Randomized Numerical Linear Algebra : A Perspective on the Field With an Eye to Software", arXiv 2023

Pseudo-code: Basic

- Compared to equations, it's more like coding.
 - ✓ Assigning operation is usual: $a \leftarrow a + 1$

- Fonts

- ✓ `\textsc{Clustering}` CLUSTERING
- ✓ `\texttt{Clustering}` Clustering

Can be used for function names

In Powerpoint, use "Consolas" font. Like this.

- Several TeX packages <https://www.overleaf.com/learn/latex/Algorithms>

```
1:  $i \leftarrow 10$ 
2: if  $i \geq 5$  then
3:    $i \leftarrow i - 1$ 
4: else
5:   if  $i \leq 3$  then
6:      $i \leftarrow i + 2$ 
7:   end if
8: end if
```

Algorithm 1 An algorithm with caption

Require: $n \geq 0$
Ensure: $y = x^n$

```
 $y \leftarrow 1$ 
 $X \leftarrow x$ 
 $N \leftarrow n$ 
while  $N \neq 0$  do
  if  $N$  is even then
     $X \leftarrow X \times X$ 
     $N \leftarrow \frac{N}{2}$  ▷ This is a comment
  else if  $N$  is odd then
     $y \leftarrow y \times X$ 
     $N \leftarrow N - 1$ 
  end if
end while
```

Algorithm 1: An algorithm with caption

Data: $n \geq 0$
Result: $y = x^n$

```
 $y \leftarrow 1$ ;
 $X \leftarrow x$ ;
 $N \leftarrow n$ ;
while  $N \neq 0$  do
  if  $N$  is even then
     $X \leftarrow X \times X$ ;
     $N \leftarrow \frac{N}{2}$ ; /* This is a comment */
  else
    if  $N$  is odd then
       $y \leftarrow y \times X$ ;
       $N \leftarrow N - 1$ ;
    end
  end
end while
end
```

Pseudo-code: Basic

- Again, consistency is important!
 - ✓ 😊 OK: $a \leftarrow a + 1 \dots b \leftarrow f(x)$
 - ✓ 😞 NO!: $a \leftarrow a + 1 \dots b = f(x)$
- You can mix a mathematical way and a coding way
 - ✓ $x \leftarrow \frac{1}{\text{Pop}(v)} \int_0^a f(\theta) d\theta$
- Don't write too much! Pseudo-code should be simple.
- Several styles
 - $v.\text{push_back}(a)$, $\text{Append}(v, a)$, $v \leftarrow [v|a]$ Assuming a row vector
 - 😞 Again, don't use italic for function: $v.\textit{push_back}(a)$

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- **What do you want to express?**
- Misc

What do you want to express?

- The design of pseudo-code depends on what you want to express.
- Consider `std::vector<int> arr` in your code
- If you use `arr` to represent a set of integers,
 - ✓ You can use a set:
 - ✓ $\mathcal{A} \leftarrow \emptyset$: initialization
 - ✓ $\mathcal{A} \leftarrow \mathcal{A} \cup \{x\}$: add
 - ✓ $\mathcal{A} \leftarrow \mathcal{A} \setminus \{y\}$: delete
- If $O(1)$ access is important, or use it in a mathematical context,
 - ✓ You should use an array (vector)
 - ✓ $\mathbf{a} \in \mathbb{R}^D$
 - ✓ $\mathbf{a}[i]$ or a_i : Implies an $O(1)$ access.

Case study: Beam search for neighbor search

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize(l)$  // remove nodes from back of  $S$  to keep its
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set \mathcal{L} containing k -approx NNs, and a set \mathcal{V} containing all the visited nodes

```
begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
     $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$ 
    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
      points to  $x_q$ 
    end if
  end while
  return [closest  $k$  points from  $\mathcal{L}; \mathcal{V}$ ]
```

DiskANN [Subramanya+, NeurIPS 19]

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

```
while has runtime budget do
   $v_i = \text{SelectNearest}(H, q)$ 
  for  $\hat{v} \in \text{Expand}(v_i, G)$  do
    if  $\hat{v} \notin V$  then
       $V := \text{Add}(V, \hat{v})$ 
       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

- All papers have totally different pseudo code for the almost same algorithm 🙄
- Hint: Explicitly state the data structure or not

Case study: Beam search for neighbor search

Algorithm 1 Search-on-Graph(G, p, q, l)
Require: graph G , start node p , query point q , candidate pool size l
Ensure: k nearest neighbors of q
1: $i = 0$, candidate pool $S = \emptyset$
2: $S.add(p)$
3: **while** $i < l$ **do**
4: $i =$ the id of the first unchecked node p_i in S
5: mark p_i as checked
6: **for all** neighbor n of p_i in G **do**
7: if n has not been visited **then**
8: $S.add(n)$
9: **end if**
10: **end for**
11: S sort in ascending order of the distance to q
12: **if** $S.size() > l$ **then**
13: $S.resize(l)$ // remove nodes from back of S to keep its size no larger than l
14: **end if**
15: **end while**
17: return the first k nodes in S

NSG [Cong+, VLDB 19]

➤ Sort the array explicitly

When need to sort, say "closest L points"

almost same algorithm 😞

➤ Hint: Explicitly state the data structure or not

Candidates are stored in a set

Algorithm 1 GreedySearch(s, x_q, k, L)
Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$
Result: Result set \mathcal{L} containing k -approx NNs, and a set \mathcal{V} containing all the visited nodes
begin
initialize sets $\mathcal{L} \leftarrow \{s\}$ and $\mathcal{V} \leftarrow \emptyset$
while $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$ **do**
let $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$
update $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$ and $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$
if $|\mathcal{L}| > L$ **then**
update \mathcal{L} to retain closest L points to x_q
end while
return [closest k points from \mathcal{L} ; \mathcal{V}]

DiskANN [Subramanya+, NeurIPS 19]

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

while has runtime budget **do**

$v_i = \text{SelectNearest}(H, q)$

for $\hat{v} \in \text{Expand}(v_i, G)$ **do**

if $\hat{v} \notin V$ **then**

$V := \text{Add}(V, \hat{v})$

$H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$

end

end

end

return TopK(V, q, k)

Learning to route [Baranchuk+, ICML 19]

Candidates are stored in a heap; automatically sorted

Case study: Beam search for neighbor search

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize()$  // remove nodes from back of  $S$  to keep its
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set \mathcal{L} containing k -approx NNs, and a set \mathcal{V} containing all the visited nodes

```
begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
            $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$ 
    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
      points to  $x_q$ 
    end if
  end while
  return [closest  $k$  points from  $\mathcal{L}; \mathcal{V}$ ]
```

DiskANN [Subramanya+, NeurIPS 19]

Checked items are stored in a set

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

```
while has runtime budget do
   $v_i = \text{SelectNearest}(H, q)$ 
  for  $\hat{v} \in \text{Expand}(v_i, G)$  do
    if  $\hat{v} \notin V$  then
       $V := \text{Add}(V, \hat{v})$ 
       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

➤ Just "check" have totally different pseudo code for the almost same algorithm 😞

➤ Hint: Explicitly state the data structure or not

Case study: Beam search for neighbor search

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize(l)$  // remove nodes from back of  $S$  to keep its
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set \mathcal{L} containing k -approx NNs, and a set \mathcal{V} containing all the visited nodes

```
begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
     $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$ 
    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
      points to  $x_q$ 
    end if
  end while
  return [closest  $k$  points from  $\mathcal{L}$ ;  $\mathcal{V}$ ]
```

DiskANN [Subramanya+, NeurIPS 19]

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

```
while has runtime budget do
   $v_i = \text{SelectNearest}(H, q)$ 
  for  $\hat{v} \in \text{Expand}(v_i, G)$  do
    if  $\hat{v} \notin V$  then
       $V := \text{Add}(V, \hat{v})$ 
       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

Visited item are simply said to be "visited"; implying an additional hidden data structure (array)

Visited items are stored in a set

almost same algorithm

➤ Hint: Explicitly state the data structure or not

Case study: Beam search for neighbor search

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize(l)$  // remove nodes from back of  $S$ 
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set \mathcal{L} containing k -approx NNs, and a set \mathcal{V} containing all the visited nodes

begin

```
initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
  let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
  update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
   $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$ 
  if  $|\mathcal{L}| > L$  then
    update  $\mathcal{L}$  to retain closest  $L$ 
```

Termination condition??

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

while has runtime budget **do**

```
 $v_i = \text{SelectNearest}(H, q)$ 
for  $\hat{v} \in \text{Expand}(v_i, G)$  do
  if  $\hat{v} \notin V$  then
     $V := \text{Add}(V, \hat{v})$ 
     $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
  end
end
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

- All papers have totally different pseudo code for the almost same algorithm 😞
- Hint: Explicitly state the data structure or not

Case study

function CSGVERTICES

Input: \mathcal{V} , \mathcal{F} : set of vertices and facets of input polyhedra

Output: \mathcal{V}_f : corresponding set of final output vertices

for F_1 **in** \mathcal{F} **do**

for v **in** F_1 **do**

if ISFINAL1(v) **then** $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v\}$

end for

for F_2 **in** \mathcal{F} **do**

$v_1, v_2 \leftarrow$ INTERSECT2FACETS(F_1, F_2)

if $\{v_1, v_2\} = \emptyset$ **then** **continue** F_2 **loop**

if ISFINAL2(v_1) **then** $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v_1\}$

if ISFINAL2(v_2) **then** $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v_2\}$

for F_3 **in** \mathcal{F} **do**

$v \leftarrow$ INTERSECTSEGMENTFACET(v_1, v_2, F_3)

if $v \neq \emptyset$ **and** ISFINAL3(v) **then** $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v\}$

end for

end for

end for

end function

- ▷ Enumerate input faces
- ▷ Order-1 candidates

- ▷ Order-2 candidates
- ▷ No intersection

- ▷ Order-3 candidates

Case study

Enumeration for a set

```
function CSGVERTICES
  Input:  $\mathcal{V}, \mathcal{F}$ : set of vertices and facets of input polyhedra
  Output:  $\mathcal{V}_f$ : corresponding set of final output vertices
  for  $F_1$  in  $\mathcal{F}$  do
    for  $v$  in  $F_1$  do
      if ISFINAL1( $v$ ) then  $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v\}$ 
    end for
    for  $F_2$  in  $\mathcal{F}$  do
       $v_1, v_2 \leftarrow \text{INTERSECT2FACETS}(F_1, F_2)$ 
      if  $\{v_1, v_2\} = \emptyset$  then continue  $F_2$  loop
      if ISFINAL2( $v_1$ ) then  $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v_1\}$ 
      if ISFINAL2( $v_2$ ) then  $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v_2\}$ 
      for  $F_3$  in  $\mathcal{F}$  do
         $v \leftarrow \text{INTERSECTSEGMENTFACET}(v_1, v_2, F_3)$ 
        if  $v \neq \emptyset$  and ISFINAL3( $v$ ) then  $\mathcal{V}_f \leftarrow \mathcal{V}_f \cup \{v\}$ 
      end for
    end for
  end for
end function
```

Set update

\textsc for a function name

Returning two variables
(This is often strange in a usual math,
but it's ok in pseudo-codes)

Comments

- ▷ Enumerate input faces
- ▷ Order-1 candidates

- ▷ Order-2 candidates
- ▷ No intersection

- ▷ Order-3 candidates

OOP style?

Algorithm 2: Lookup(\mathcal{T}, k)

Input: \mathcal{T} : the LIPP index, k : the lookup key

Output: *isFound*: indicates whether k is found,
 e : the entry containing k if found

```
1 begin
2    $n \leftarrow$  the root node of  $\mathcal{T}$ ;
3    $e \leftarrow n.\mathcal{E}[n.\mathcal{M}(k)]$ ;
4   while  $type(e) == NODE$  do
5      $n \leftarrow$  the node pointed to from entry  $e$ ;
6      $e \leftarrow n.\mathcal{E}[n.\mathcal{M}(k)]$ ;
7   if  $type(e) == DATA$  then
8      $k' \leftarrow$  the key in entry  $e$ ;
9     if  $k == k'$  then
10      return  $\langle True, e \rangle$ ;
11  return  $\langle False, e \rangle$ ;
12 end
```

- It's ok to define a class as well
- Member variables can be accessed via a usual "." notation.
- But don't make the thing complex.

Equations

- Basic notation for variables
- String
- Tips for sets
- Inputs/outputs of functions
- Subscript/superscript
- Don't write numpy
- Misc

Pseudo codes

- Basic
- What do you want to express?
- **Misc**

Misc

- Consider what are global variables!
- You can use a sentence to describe a difficult concept, e.g.,
 - $\mathcal{J} \leftarrow$ Identifiers of top K smallest values of \mathbf{x}
 - $\mathbf{T} \leftarrow (N_1 \times N_2 \times N_3)$ empty arrays of B-Trees.
- Intentionally, you can use **almost-code** style.

Algorithm 1 SimSiam Pseudocode, PyTorch-like

```
# f: backbone + projection mlp
# h: prediction mlp

for x in loader: # load a minibatch x with n samples
    x1, x2 = aug(x), aug(x) # random augmentation
    z1, z2 = f(x1), f(x2) # projections, n-by-d
    p1, p2 = h(z1), h(z2) # predictions, n-by-d

    L = D(p1, z2)/2 + D(p2, z1)/2 # loss

    L.backward() # back-propagate
    update(f, h) # SGD update

def D(p, z): # negative cosine similarity
    z = z.detach() # stop gradient

    p = normalize(p, dim=1) # l2-normalize
    z = normalize(z, dim=1) # l2-normalize
    return -(p*z).sum(dim=1).mean()
```

- Their contribution is it's easy to implement the algorithm in PyTorch
- Complicated idea, such as "detach", can be explained in one line.
- But be careful! Can future readers understand?

Misc

- If you use almost-code style, use the minted package.
- https://www.overleaf.com/learn/latex/Code_Highlighting_with_minted

```
\documentclass{article}
\usepackage{minted}
\begin{document}
\begin{minted}{python}
import numpy as np

def incmatrix(genl1,genl2):
    m = len(genl1)
    n = len(genl2)
    M = None #to become the incidence matrix
    VT = np.zeros((n*m,1), int) #dummy variable

    #compute the bitwise xor matrix
    M1 = bitxormatrix(genl1)
    M2 = np.triu(bitxormatrix(genl2),1)

    for i in range(m-1):
        for j in range(i+1, m):
            [r,c] = np.where(M2 == M1[i,j])
            for k in range(len(r)):
                VT[(i)*n + r[k]] = 1;
                VT[(i)*n + c[k]] = 1;
                VT[(j)*n + r[k]] = 1;
                VT[(j)*n + c[k]] = 1;

            if M is None:
                M = np.copy(VT)
            else:
                M = np.concatenate((M, VT), 1)

    VT = np.zeros((n*m,1), int)
```



```
import numpy as np

def incmatrix(genl1,genl2):
    m = len(genl1)
    n = len(genl2)
    M = None #to become the incidence matrix
    VT = np.zeros((n*m,1), int) #dummy variable

    #compute the bitwise xor matrix
    M1 = bitxormatrix(genl1)
    M2 = np.triu(bitxormatrix(genl2),1)

    for i in range(m-1):
        for j in range(i+1, m):
            [r,c] = np.where(M2 == M1[i,j])
            for k in range(len(r)):
                VT[(i)*n + r[k]] = 1;
                VT[(i)*n + c[k]] = 1;
                VT[(j)*n + r[k]] = 1;
                VT[(i)*n + c[k]] = 1;
```

Schedule

| Date (2026) | Contents | Presented by |
|---------------------------|---|---|
| Week 1, Apr 15 | Introduction. Review of fundamental concepts | Yusuke, Koya, Yuki, Jun |
| Week 2, Apr 22 | Equations and pseudo-codes | Yusuke Matsui |
| Week 3, May 13 | Slides | Koya Narumi |
| Week 4, May 20 | Research community | Jun Kato |
| Week 5, May 27 | Figures | Koya Narumi |
| Week 6, June 1 | 3DCG illustrations | Yuki Koyama |
| Week 7, June 10 | Invited Talk 1: Research tips in the private sector: From topic selection to social implementation (temp.) | Dr. Antonio Tejero de Pablos (CyberAgent) |
| Week 8, June 17 | Tables and plots | Yusuke Matsui |
| Week 9, June 24 | Invited Talk 2: Portable and Reproducible Research Environments in the Age of AI Agents | Dr. Mai Nishimura (OSX) |
| Week 10, July 1 | DevOps for research | Jun Kato |
| Week 11, July 8 | Videos | Koya Narumi |
| Week 12, July 15 | Final presentations | YOU |

